

#### COSC160: Data Structures Linked Lists

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### Outline

- I. Data Structure (lower-level) Implementations
  - I. Chaining (Linked Lists)
    - I. Design Decisions and Implications on Complexity
      - I. Example: Tail Pointer
  - II. Arrays
    - I. Design Decisions and Implications on Complexity
- II. Design Decisions
  - I. Arrays vs. Chaining
- III. Interesting Applications
  - I. Polynomial Arithmetic and Evaluation



#### Linear Structures

- Linear structures are generally implemented using arrays or chaining
  - Contiguous vs. non-contiguous in memory
    - Static size vs. dynamic size
  - When designing a data structure, one must consider these low-level designs
    - One major decision: contiguous vs non-contiguous (ie array vs chain)
    - This design decision will have implications with respect to usage and efficiency
  - There are many design decisions that a programmer must consider
    - Some are standard, others may be tailor-made for the problem/solution on-hand.



# List Class Implemented by Chaining (Linked Lists)

- Design Decisions
  - Single Link, next
  - Head pointer maintained

- Standard Operations
  - Insert(item, index)
  - Remove(index)
  - Retrieve(index)



### Inserting into a Linked List

- Lets investigate three cases of the insertion operation
  - Insert to front of list
  - Insert to middle of list
  - Insert to end of list



# Inserting into front of a Linked List

- Insert to front of list
- Approximately how many computational steps?
- Does the number of computational steps depend on the size of the list?



• Time complexity?



## Inserting into middle of Linked List



Approximately how many computational steps?

Does the number of computational steps depend on the size of the list?

Time complexity?





# Inserting to end of Linked List

- Approximately how many computational steps?
- Does the number of computational steps depend on the size of the list?









#### Design Decisions

- One major goal of a data structure is to provide an efficient way to manage data.
- Can we improve the efficiency of any of these operations -- YES



### Insert to End (List with Tail Pointer)

- Significant time reduction.
- What is the time complexity now?





# *List Variations: Singly Linked vs. Doubly Linked Lists*

- Differences
  - Doubly Linked List
    - Can traverse a list in both directions at added cost of 1 pointer per node
    - Slightly simplify iteration scheme for insertions and removals



# Singly Linked vs. Doubly Linked Lists



• Doubly Linked List



# Why a double link?

- May be useful in specific situations.
  - Remove last node

- A prev pointer, may reduce retrieval and insertion times at the cost (space) of 1 extra pointer per node.
  - Create an efficient algorithm: Node\* get(int index)



# List Class Implemented as Array

- Assume the following design decisions
  - Array initialized to "large" capacity, chosen appropriately based on application
  - No gaps in the list (items shift after a removal)
  - Assumes index = 0 is beginning of list
  - Should maintain index for end (size) of list or numltems
  - Operations
    - Insert(item, index)
    - Remove(index)
    - Retrieve(index)
  - Concerns: What happens if insertion exceeds capacity?

void



### Inserting into an Array

- Lets investigate three cases of the insertion operation
  - Insert to front of array
  - Insert to middle of array
  - Insert to end of array



# Insertion to Front of Array

Add 8 to front

 Approximately how many computational steps?



 Does the number of computational steps depend on the size of the list?



• Time complexity?





# Insert to Middle of Array

- Approximately how many computational steps?
- Does the number of computational steps depend on the size of the list?





• Time complexity?





# Insert to End of Array

size: 4

- Approximately how many ٠ computational steps?
- Does the number of ٠ computational steps depend on the size of the list? Note: Assumes there is enough allocated space.









# "Worst" Case Time Complexity for a List Class

• Exercise: try to fill out table for best case and average case.

IMPLEMENTATION	CHAINING	ARRAYS
INSERT FRONT	Θ(1)	$\Theta(n)$
INSERT MIDDLE	$\Theta(n)$	$\Theta(n)$
INSERT BACK	$\Theta(1)^*$	Θ(1)**
RETRIEVE FRONT	Θ(1)	Θ(1)
RETRIEVE MIDDLE	$\Theta(n)$	Θ(1)
RETRIEVE BACK	$\Theta(1)^*$	Θ(1)
REMOVE FRONT	Θ(1)	$\Theta(n)$
REMOVE MIDDLE	$\Theta(n)$	$\Theta(n)$
REMOVE BACK	Θ(1)	Θ(1)
* Assumes tail ptr		
** Assumes alloc		



# Basic Lists: Array vs Linked List Implementation

- Example: Design a list class for use to store information for 250 tenants in apt. building.
- Design an aptList Class with the following operations
  - Insert(item, location): bool
  - Remove(item): bool
  - Find(item): int
- How would you implement this structure as a class?
  - Array or Chain? Why?
  - What member variables would you include? Why?
- We could keep the data in the list in some order ...

